



BGP Configuration

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COS 561: Advanced Computer Networks

<http://www.cs.princeton.edu/courses/archive/fall10/cos561/>

Local Control vs. Global Properties



The Internet is a “network of networks”

- ~35,000 separately administered networks
- Competitive cooperation for e2e reachability

Local Control

Intradomain routing,
interdomain policies

Global Properties

Performance, security,
reliability, scalability





Two-Tiered Routing Architecture

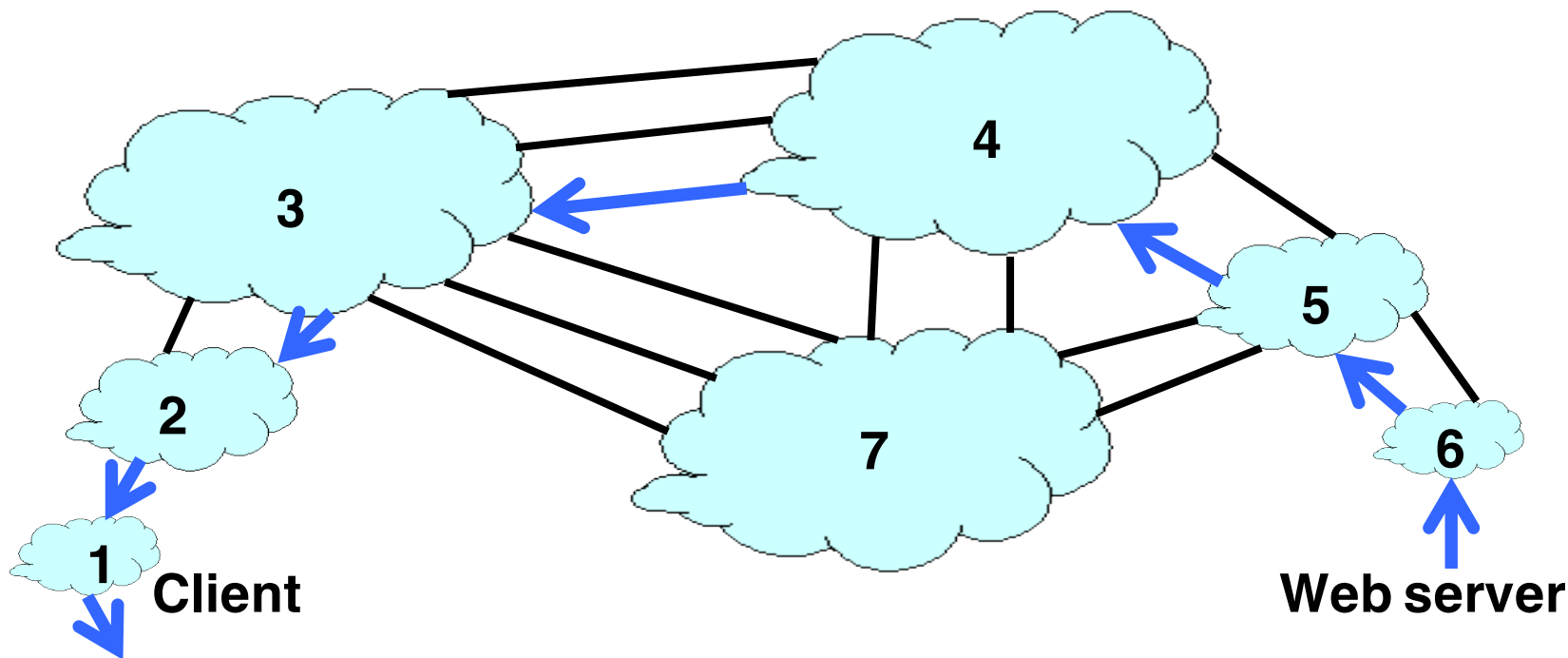
- Goal: distributed management of resources
 - Internetworking of multiple networks
 - Networks under separate administrative control
- *Intradomain*: inside a region of control
 - Routers configured to achieve a common goal
 - Okay for routers to share *topology* information
 - Different ASes can run different protocols
- *Interdomain*: between regions of control
 - ASes have different (maybe conflicting) goals
 - Routers only share *reachability* information



Internet Structure

Autonomous Systems (ASes)

- AS-level topology
 - Nodes are Autonomous Systems (ASes)
 - Destinations are prefixes (e.g., 12.0.0.0/8)
 - Edges are links and business relationships





AS Numbers (ASNs)

**ASNs are 16 bit values (or 32-bit).
64512 through 65535 are “private”**

Currently around 35,000 in use.

- **Level 3: 1**
- **MIT: 3**
- **Harvard: 11**
- **Yale: 29**
- **Princeton: 88**
- **AT&T: 7018, 6341, 5074, ...**
- **Verizon: 701, 702, 284, 12199, ...**
- **Sprint: 1239, 1240, 6211, 6242, ...**
- **...**

Business Relationships Between ASes



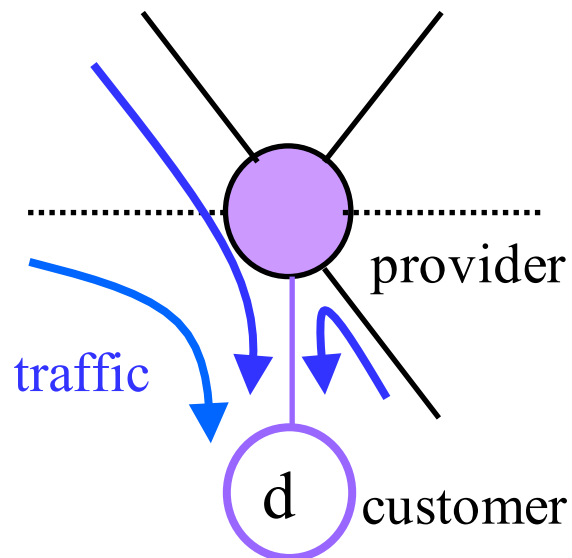
- Neighboring ASes have business contracts
 - How much traffic to carry
 - Which destinations to reach
 - How much money to pay
- Common business relationships
 - Customer-provider
 - Peer-peer
 - Backup
 - Sibling

Customer-Provider Relationship

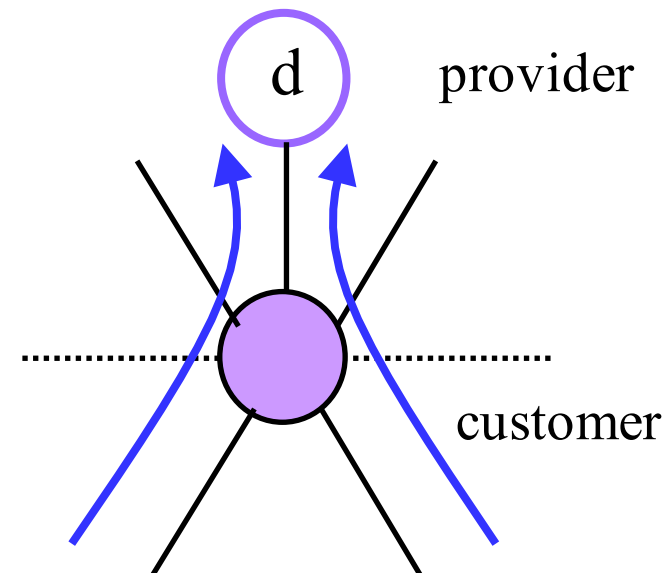


- Customer needs to be reachable from everyone
 - Provider ensures all neighbors can reach the customer
- Customer does not want to provide transit service
 - Customer does not let its providers send traffic through it

Traffic **to** the customer



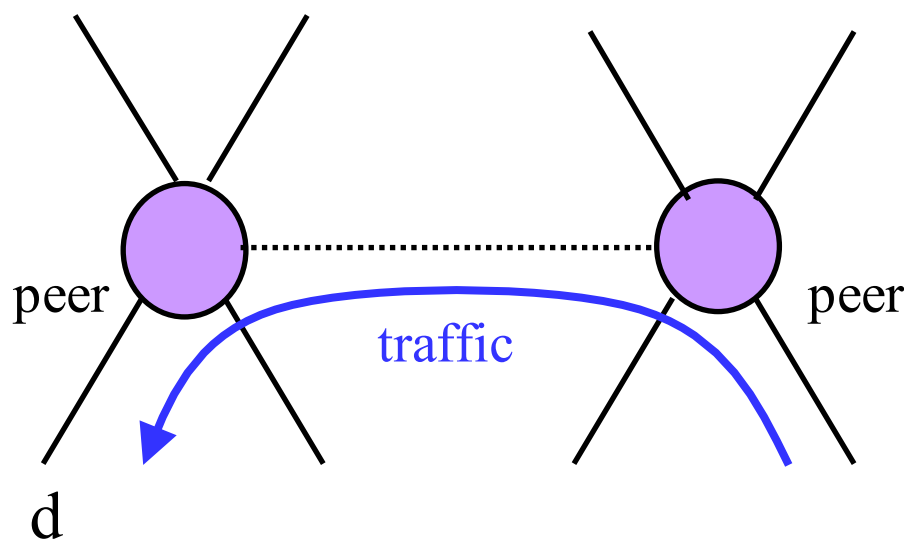
Traffic **from** the customer



Peer-Peer Relationship

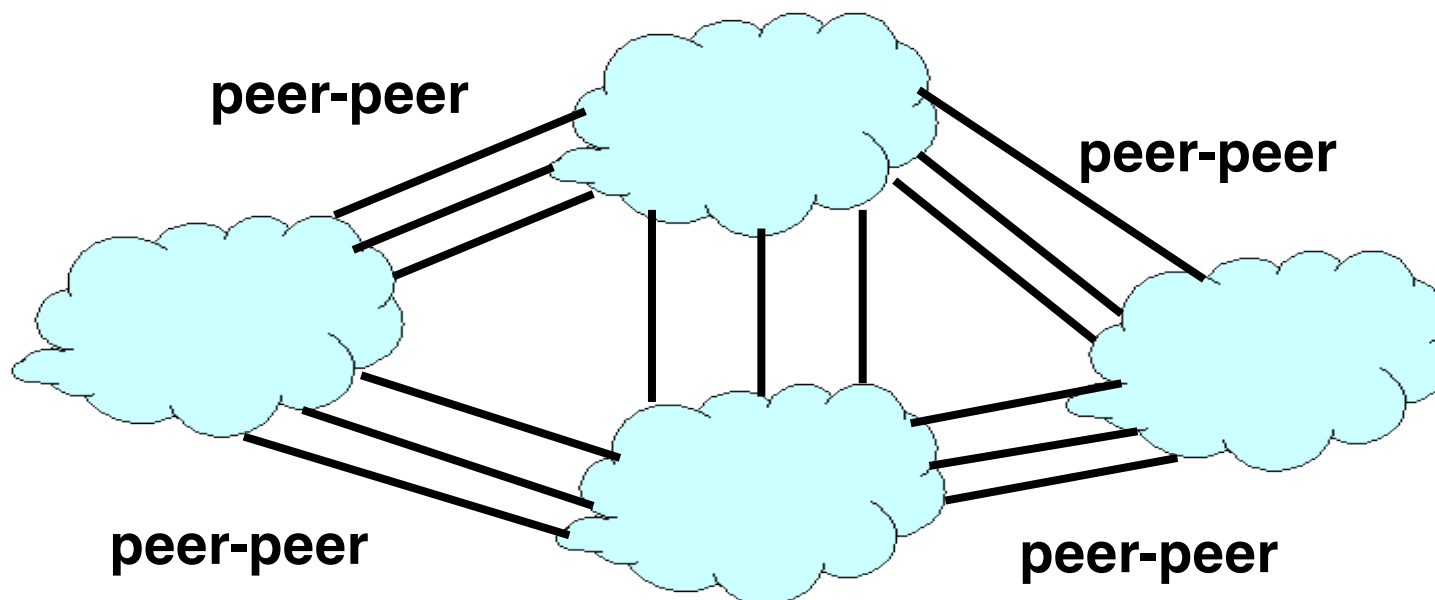
- Peers exchange traffic between customers
 - AS lets its peer reach (only) its customers
 - AS can reach its peer's customers
 - Often the relationship is settlement-free (i.e., no \$\$\$)

Traffic to/from the peer and its customers



AS Structure: Tier-1 Providers

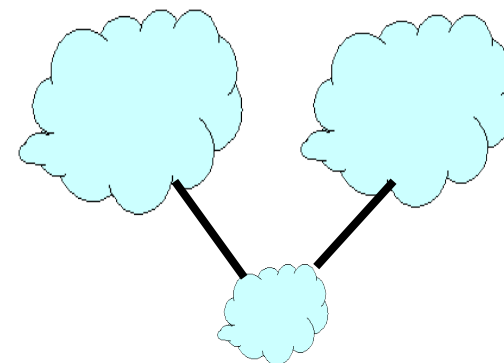
- Top of the Internet hierarchy
 - Has no upstream provider of its own
 - Typically has a large (inter)national backbone
 - Around 10-12 ASes: AT&T, Sprint, Level 3, ...





AS Structure: Other ASes

- Lower-layer providers (tier-2, ...)
 - Provide transit service to downstream customers
 - But need at least one provider of their own
 - Typically have national or regional scope
 - E.g., Minnesota Regional Network
 - Includes a few thousand ASes
- Stub ASes
 - Do not provide transit service
 - Connect to upstream provider(s)
 - Most ASes (e.g., 85-90%)
 - E.g., Princeton



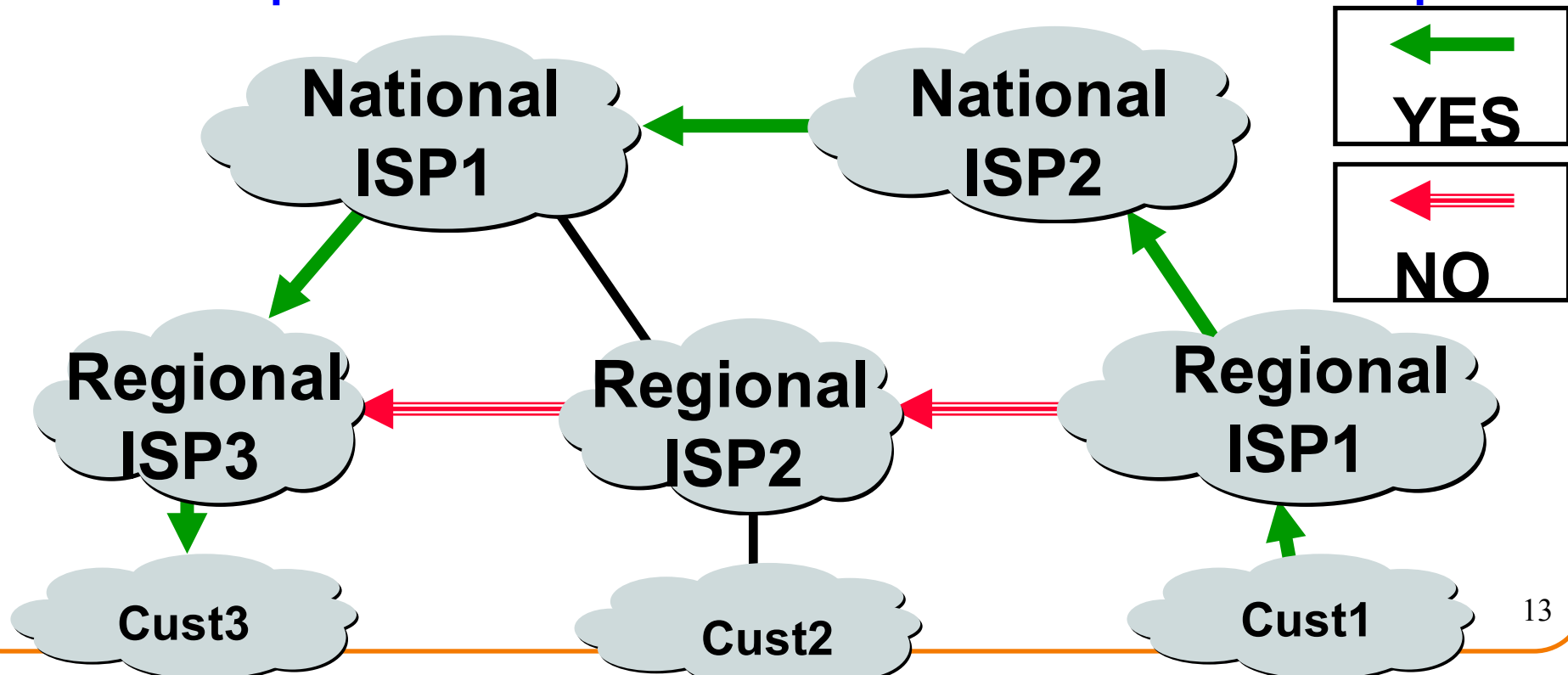


Policy-Based Path-Vector Routing

Shortest-Path Routing is Restrictive

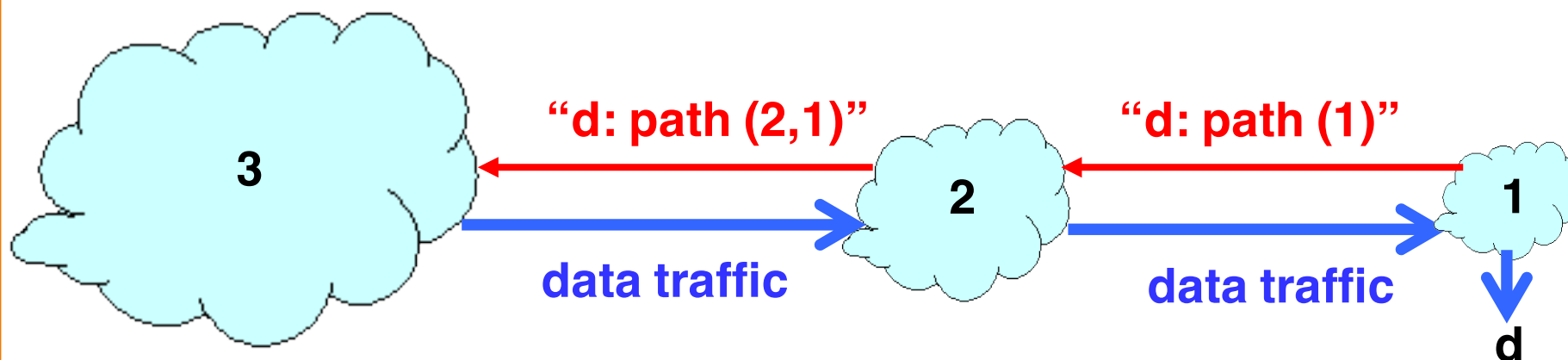


- All traffic must travel on shortest paths
- All nodes need common notion of link costs
- Incompatible with commercial relationships



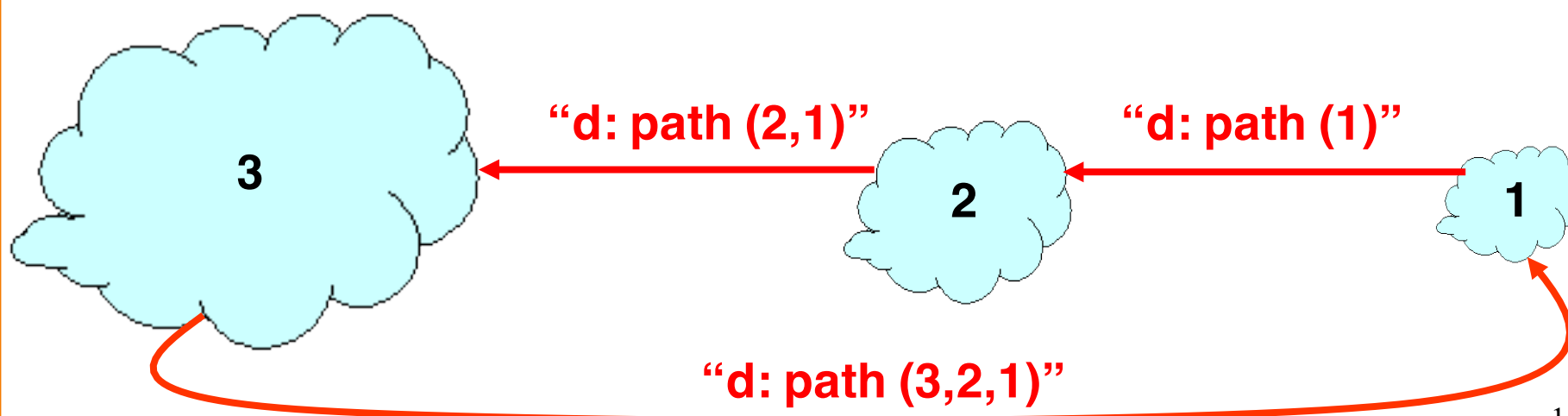
Path-Vector Routing

- Extension of distance-vector routing
 - Support flexible routing policies
 - Faster convergence (avoid count-to-infinity)
- Key idea: advertise the entire path
 - Distance vector: send *distance metric* per dest d
 - Path vector: send the *entire path* for each dest d



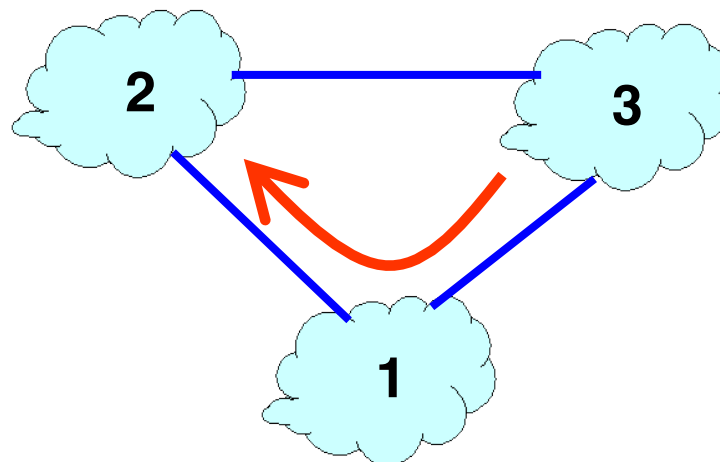
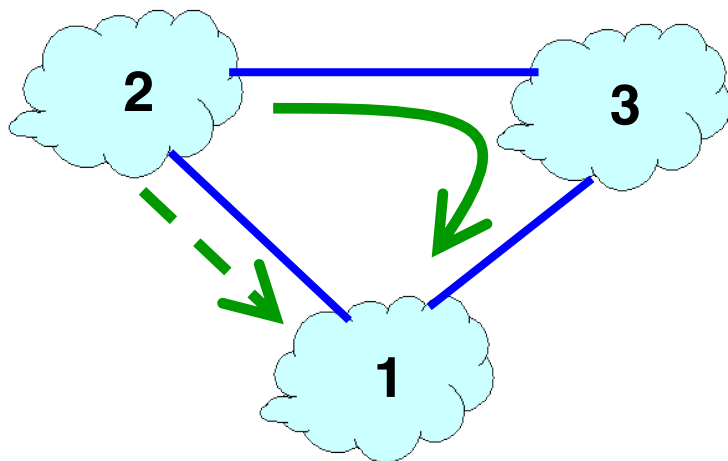
Faster Loop Detection

- Node can easily detect a loop
 - Look for its own node identifier in the path
 - E.g., node 1 sees itself in the path “3, 2, 1”
- Node can simply discard paths with loops
 - E.g., node 1 simply discards the advertisement



Flexible Policies

- Each node can apply local policies
 - Path selection: Which path to use?
 - Path export: Whether to advertise the path?
- Examples
 - Node 2 may prefer the path “2, 3, 1” over “2, 1”
 - Node 1 may not let node 3 hear the path “1, 2”





Border Gateway Protocol



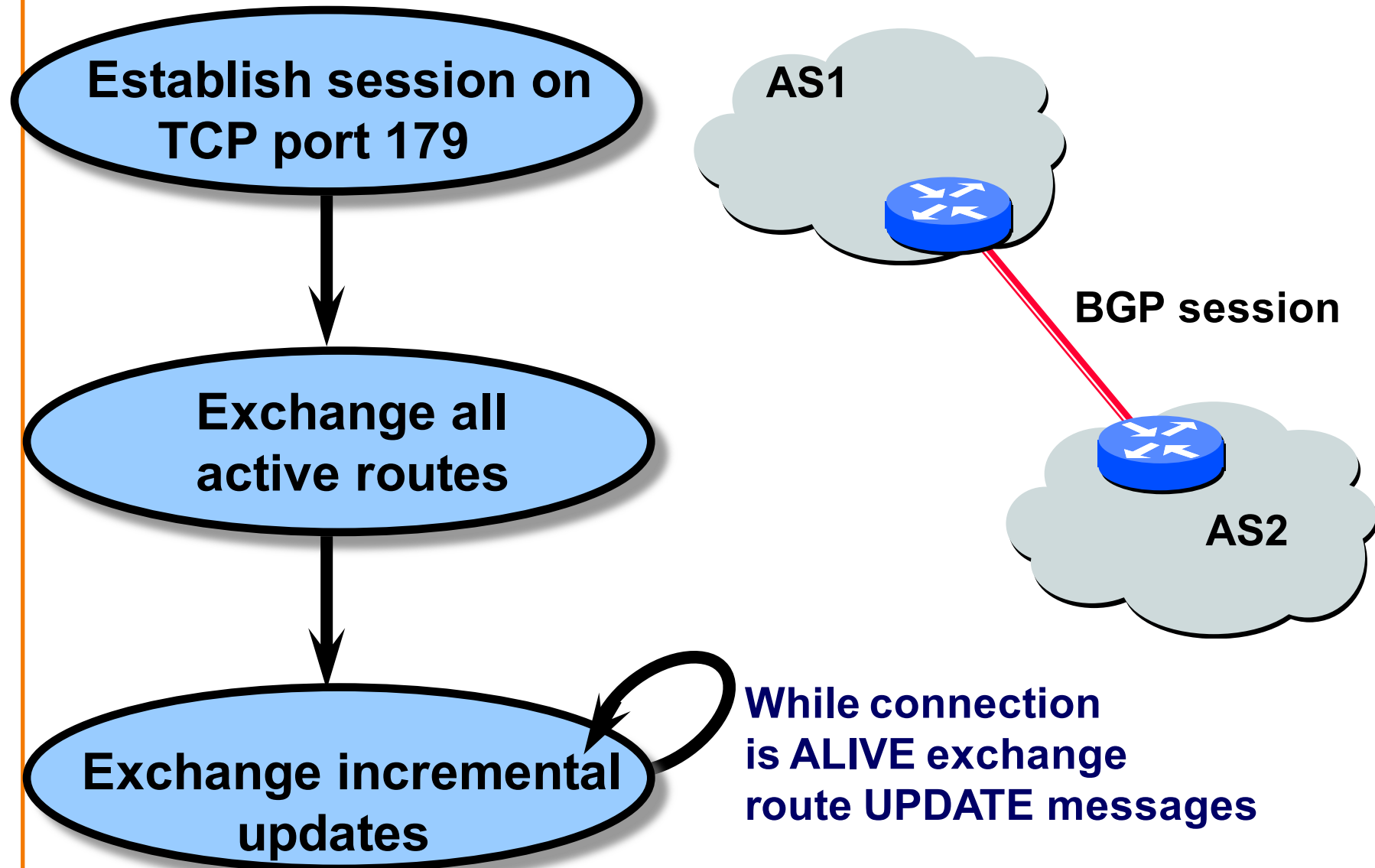
Border Gateway Protocol

- Prefix-based path-vector protocol
- Policy-based routing based on AS Paths
- Evolved during the past 20 years

- **1989 : BGP-1 [RFC 1105], replacement for EGP**
- **1990 : BGP-2 [RFC 1163]**
- **1991 : BGP-3 [RFC 1267]**
- **1995 : BGP-4 [RFC 1771], support for CIDR**
- **2006 : BGP-4 [RFC 4271], update**

“BGP at 18”: <http://www.youtube.com/watch?v=HAOVNYSnL7k> 18

BGP Operations





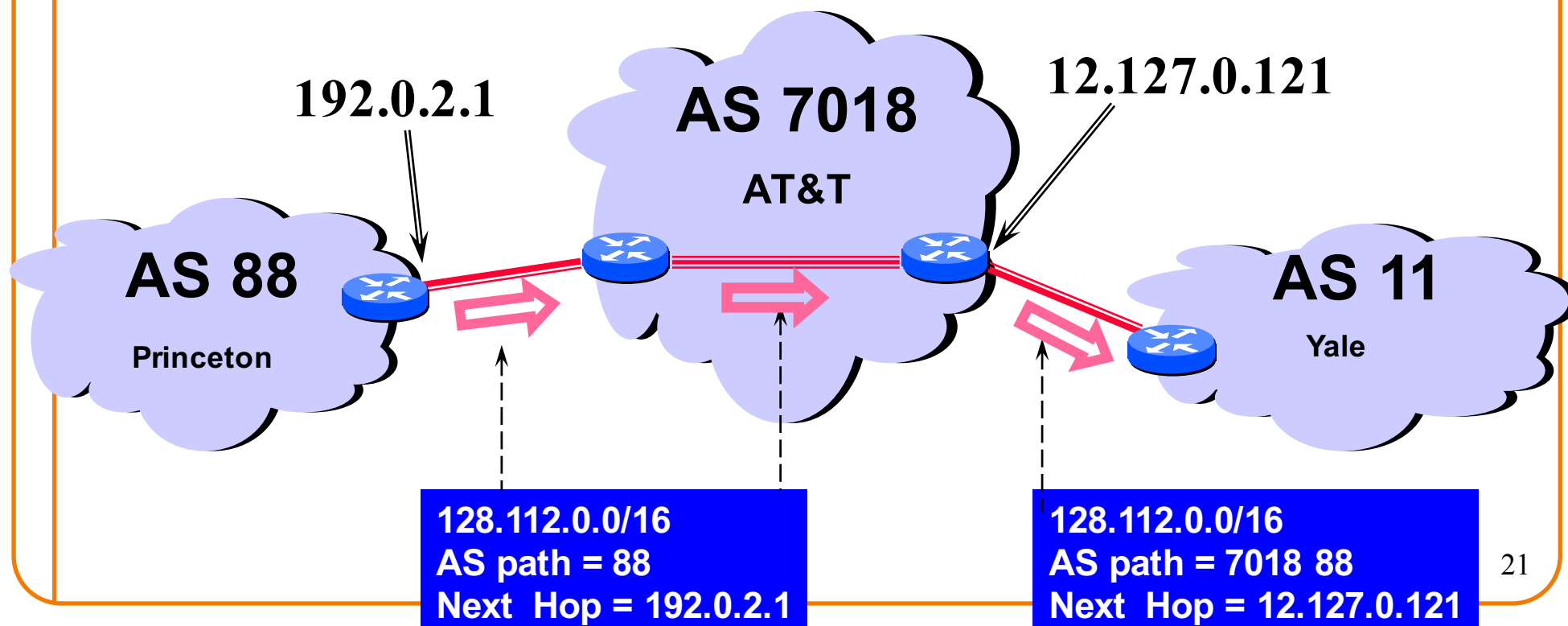
Incremental Protocol

- A node learns multiple paths to destination
 - Stores all of the routes in a routing table
 - Applies policy to select a single active route
 - ... and may advertise the route to its neighbors
- Incremental updates
 - Announcement
 - Upon selecting a new active route, add node id to path
 - ... and (optionally) advertise to each neighbor
 - Withdrawal
 - If the active route is no longer available
 - ... send a withdrawal message to the neighbors



BGP Route

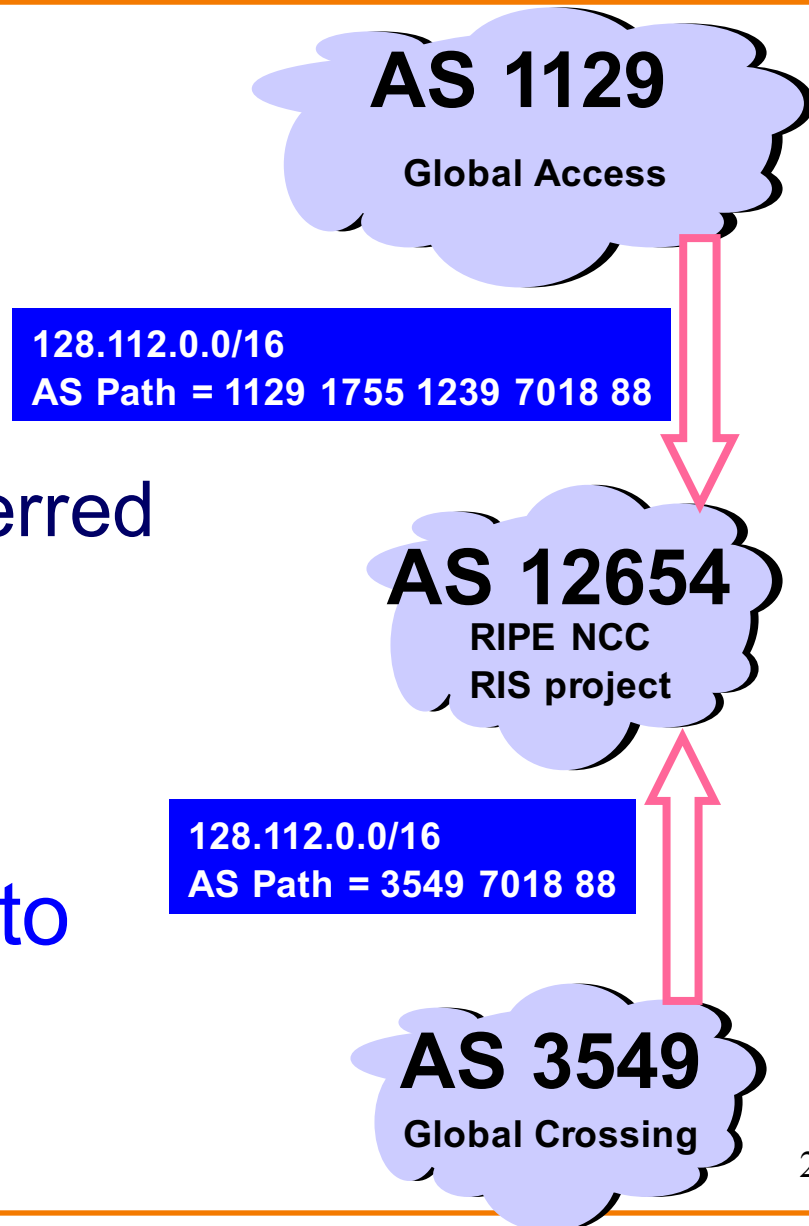
- Destination prefix (e.g., 128.112.0.0/16)
- Route attributes, including
 - AS path (e.g., “7018 88”)
 - Next-hop IP address (e.g., 12.127.0.121)



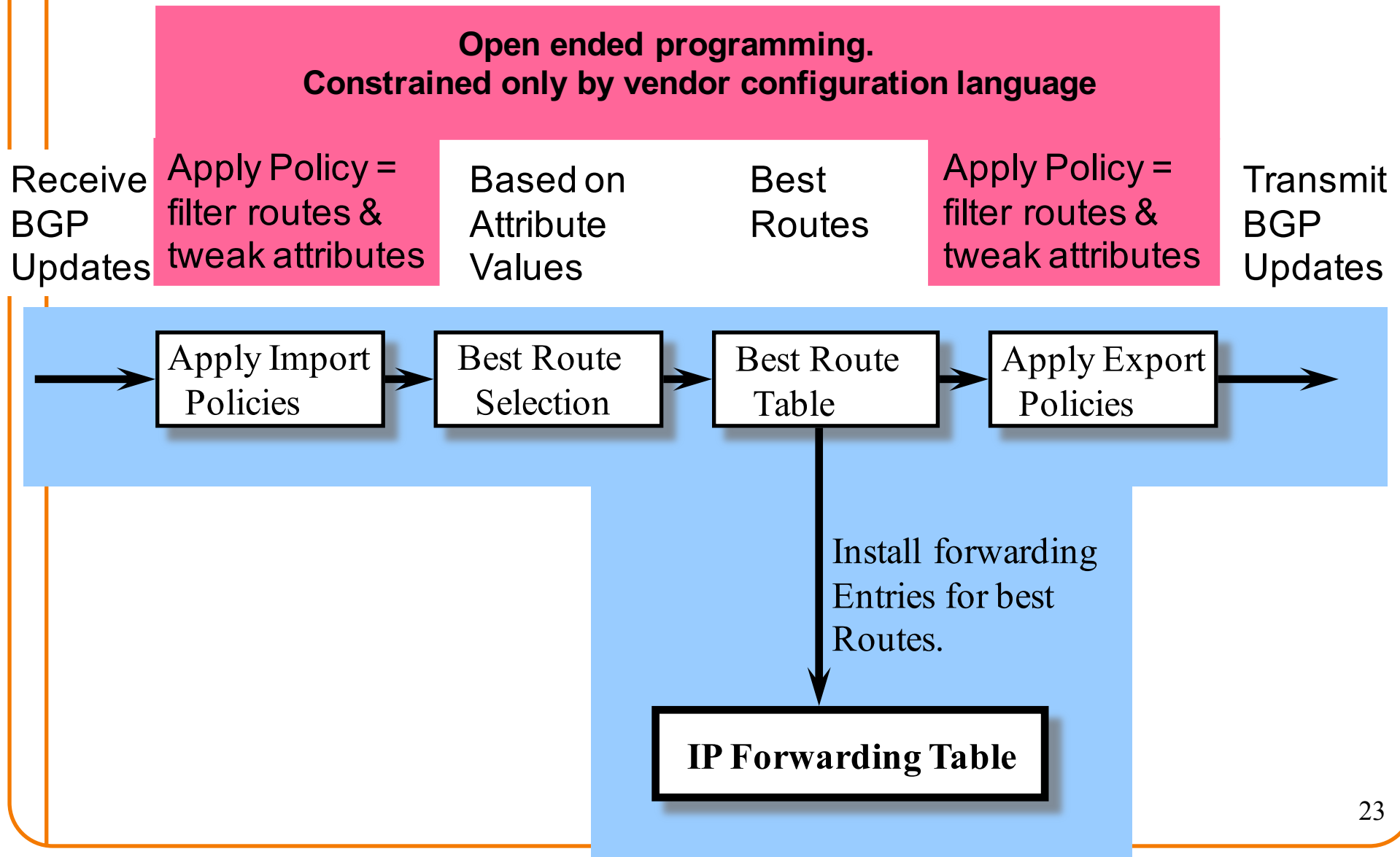
BGP Path Selection



- Simplest case
 - Shortest AS path
 - Arbitrary tie break
- Example
 - Three-hop AS path preferred over a five-hop AS path
 - AS 12654 prefers path through Global Crossing
- But, BGP is not limited to shortest-path routing
 - Policy-based routing



BGP Policy: Influencing Decisions



BGP Policy: Applying Policy to Routes



- Import policy

- Filter unwanted routes from neighbor

- E.g. prefix that your customer doesn't own

- Manipulate attributes to influence path selection

- E.g., assign local preference to favored routes

- Export policy

- Filter routes you don't want to tell your neighbor

- E.g., don't tell a peer a route learned from other peer

- Manipulate attributes to control what they see

- E.g., make a path look artificially longer than it is

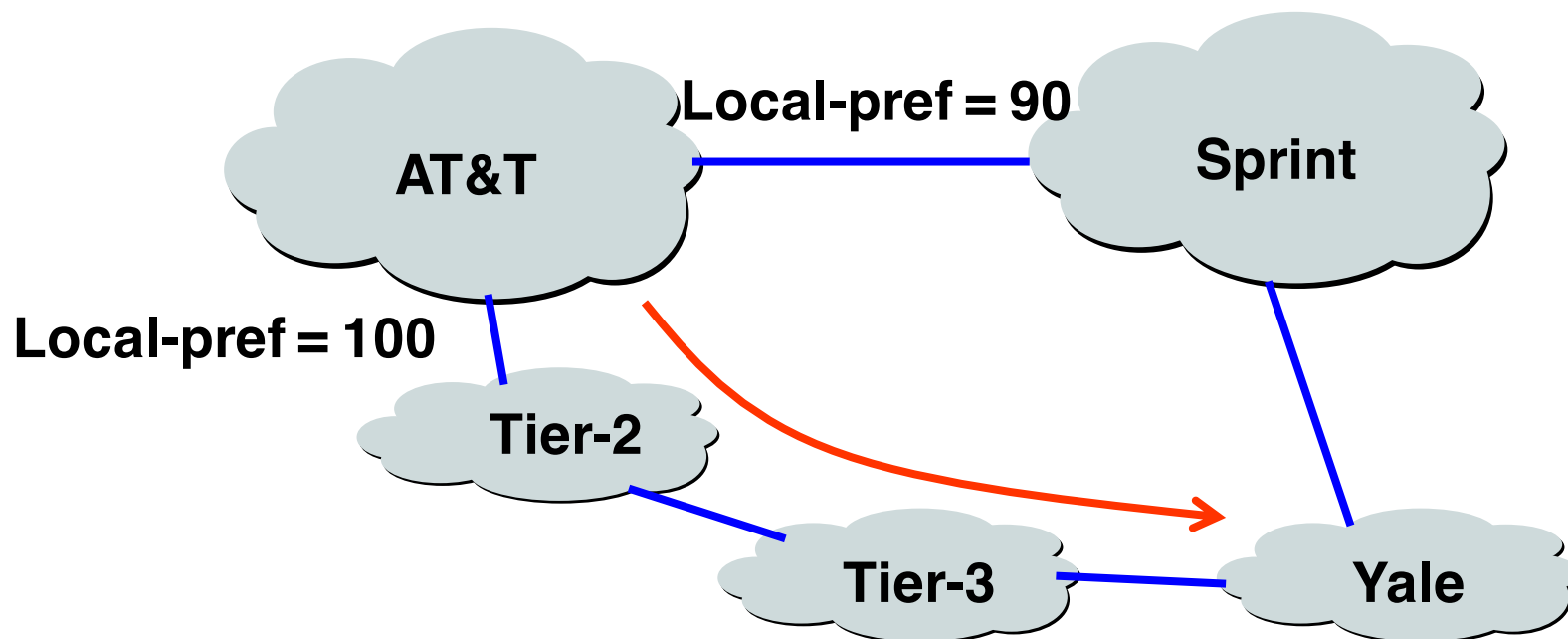


BGP Policy Examples



Import Policy: Local Preference

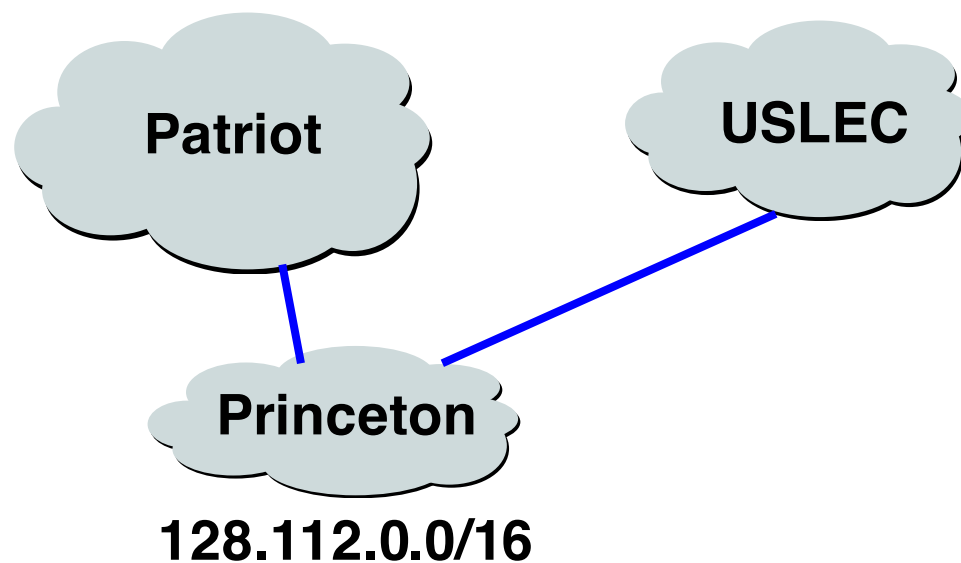
- Favor one path over another
 - Override the influence of AS path length
 - Apply local policies to prefer a path
- Example: prefer customer over peer





Import Policy: Filtering

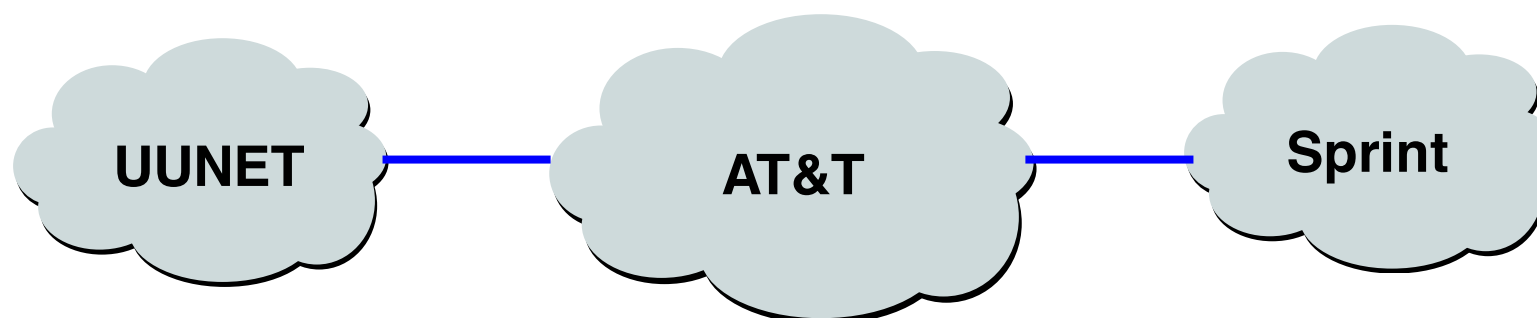
- Discard some route announcements
 - Detect configuration mistakes and attacks
- Examples on session to a customer
 - Discard route if customer doesn't own the prefix
 - Discard route containing other large ISPs





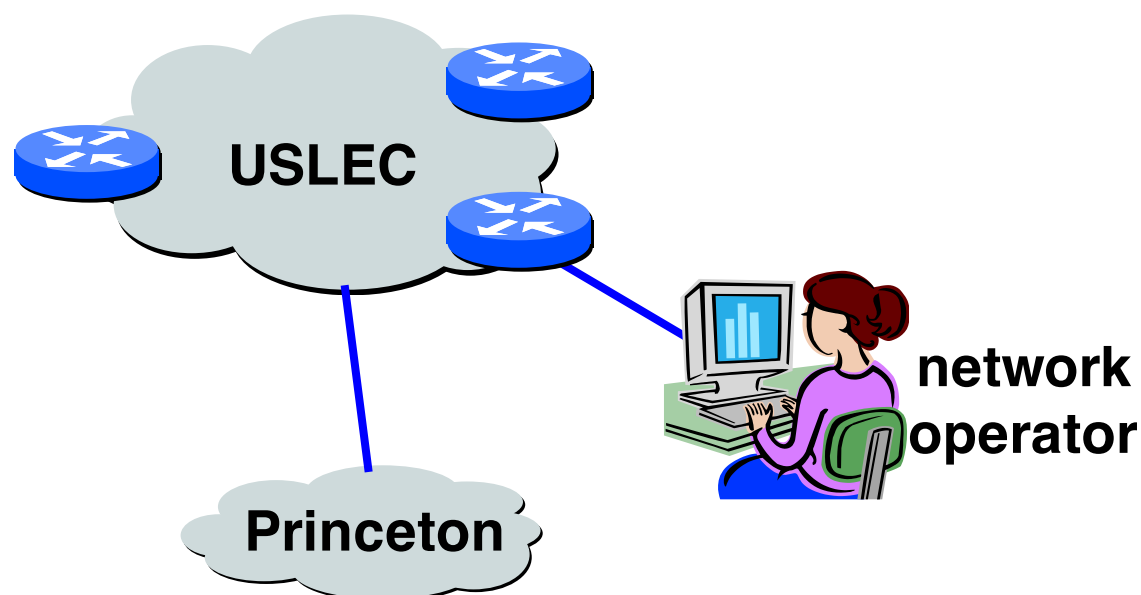
Export Policy: Filtering

- Discard some route announcements
 - Limit propagation of routing information
- Examples
 - Don't announce routes from one peer to another



Export Policy: Filtering

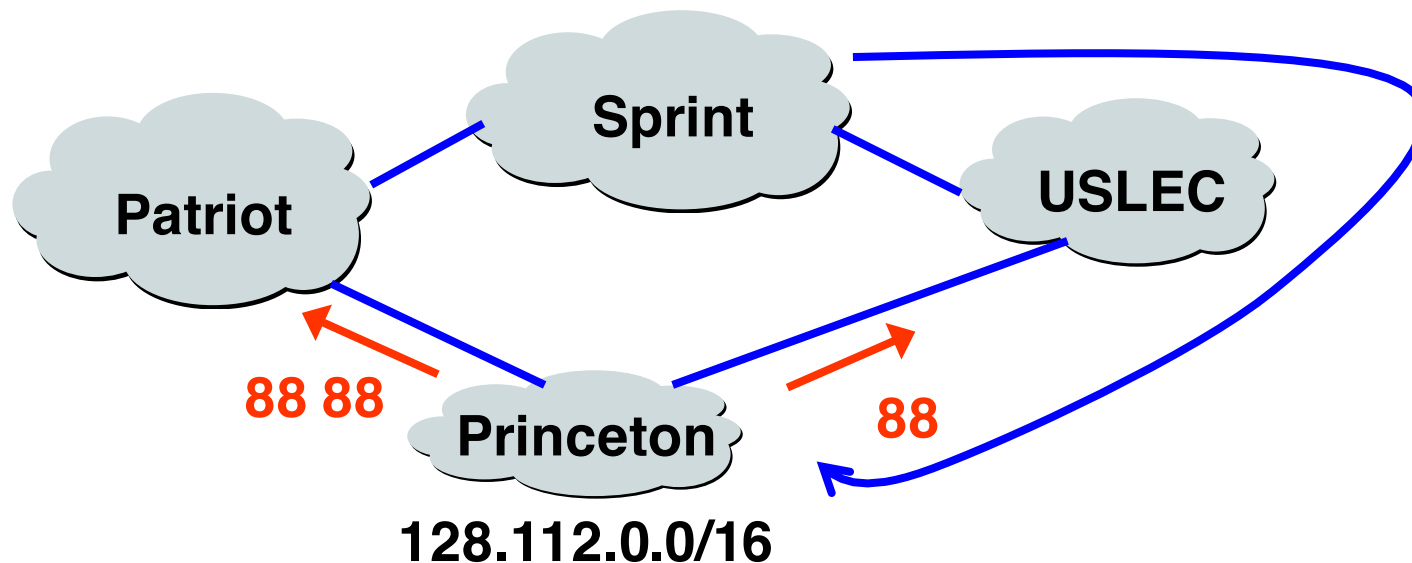
- Discard some route announcements
 - Limit propagation of routing information
- Examples
 - Don't announce routes for network-management hosts or the underlying routers themselves



Export Policy: Attribute Manipulation



- Modify attributes of the active route
 - To influence the way other ASes behave
- Example: AS prepending
 - Artificially inflate AS path length seen by others
 - Convince some ASes to send traffic another way





BGP Policy Configuration

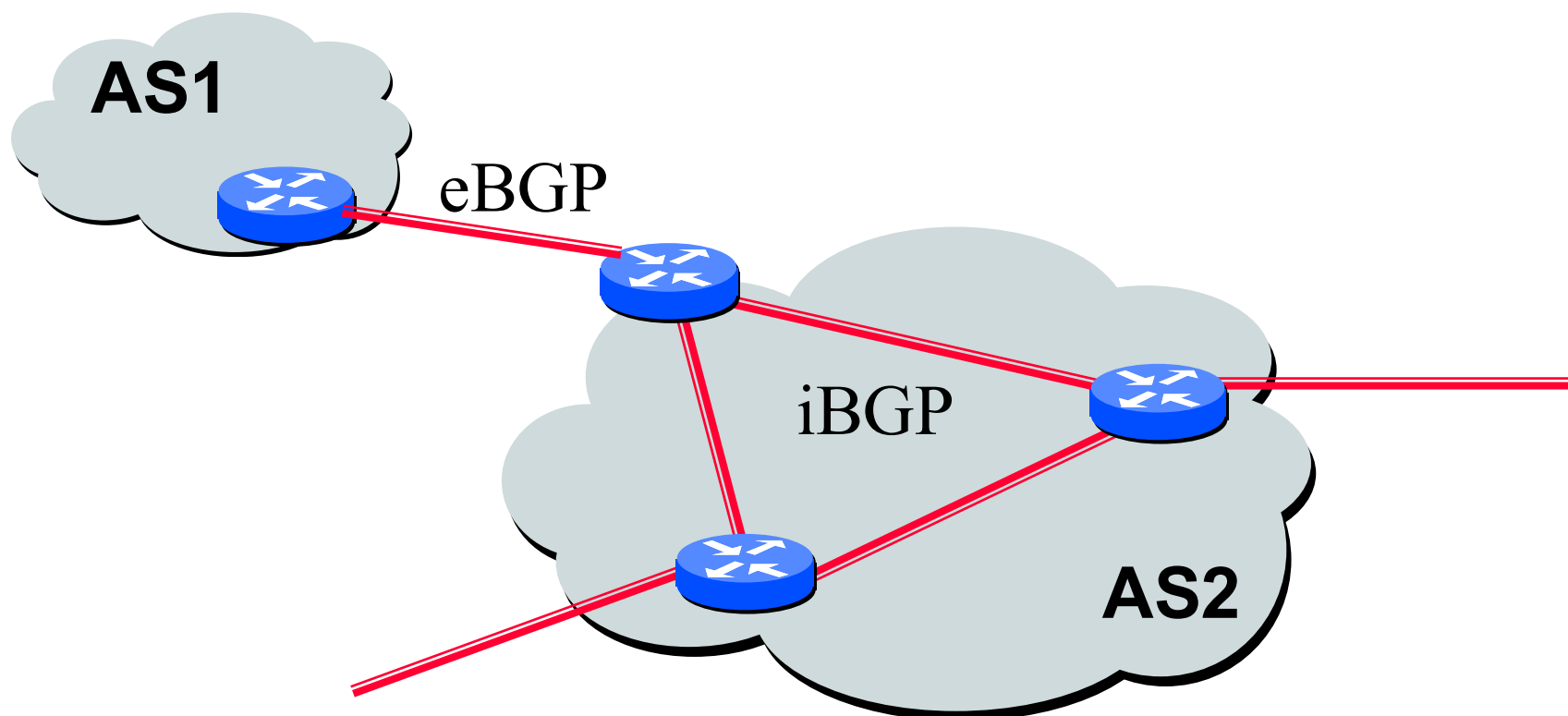
- Policy languages are vendor-specific
 - Not part of the BGP protocol specification
 - Different languages for Cisco, Juniper, etc.
- Still, all languages have some key features
 - Policy as a list of clauses
 - Each clause matches on route attributes
 - ... and discards or modifies the matching routes
- Configuration done by human operators
 - Implementing the policies of their AS
 - Biz relationships, traffic engineering, security, ..



BGP Inside an AS

An AS is Not a Single Node

- Multiple routers in an AS
 - Need to distribute BGP information within the AS
 - Internal BGP (iBGP) sessions between routers



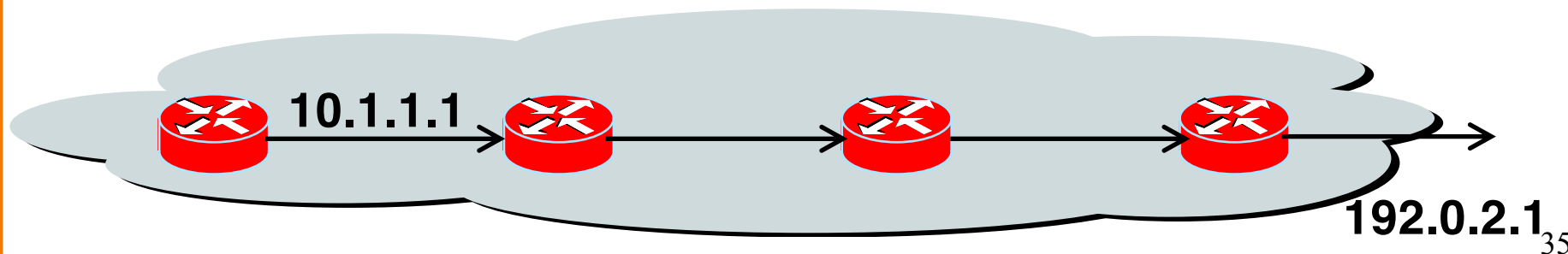
- Both routers prefer the path through AS 100 on the left
- ... even though the right router learns an external path





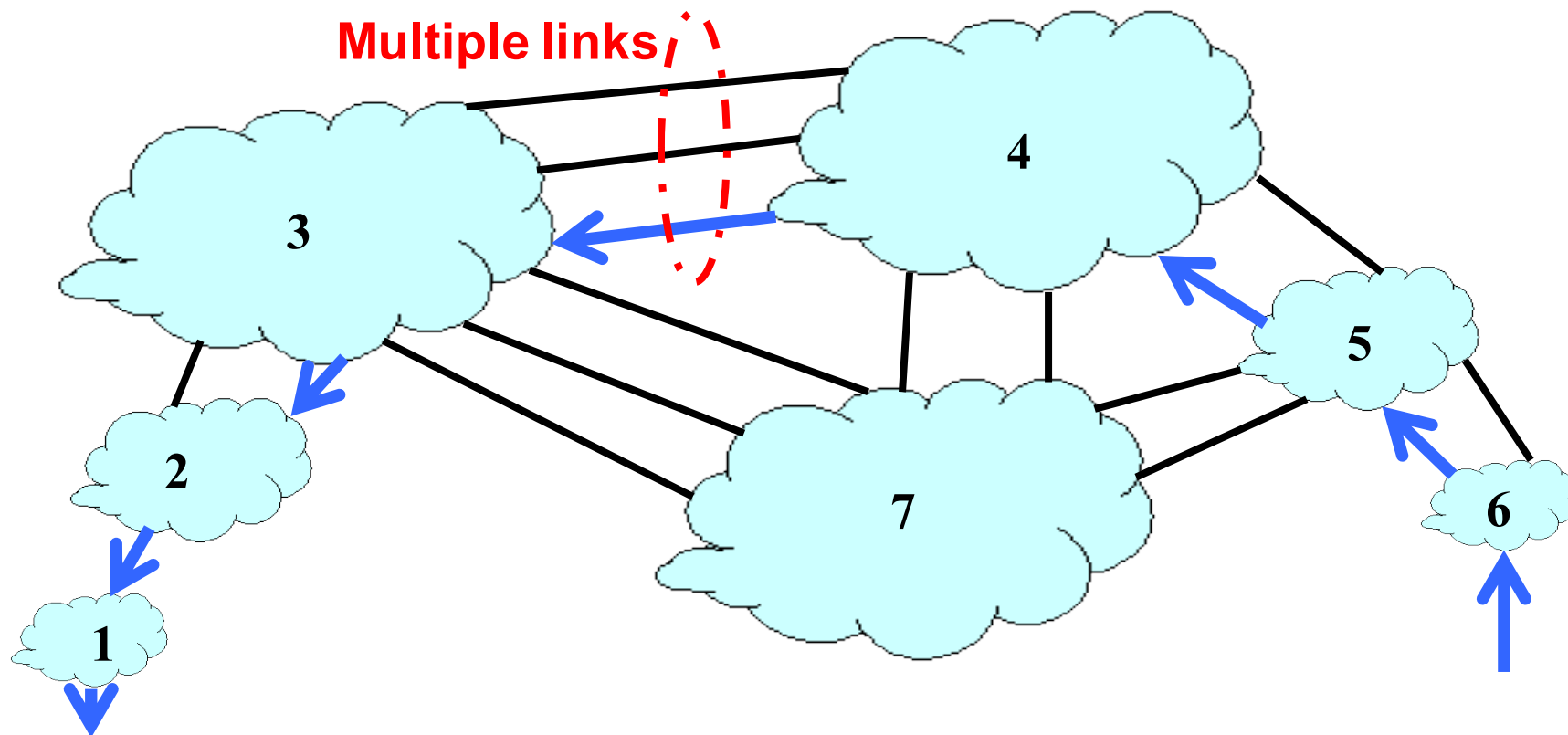
Joining BGP and IGP Information

- Border Gateway Protocol (BGP)
 - Announces reachability to external destinations
 - Maps a destination prefix to an egress point
 - 128.112.0.0/16 reached via 192.0.2.1
- Interior Gateway Protocol (IGP)
 - Used to compute paths within the AS
 - Maps an egress point to an outgoing link
 - 192.0.2.1 reached via 10.1.1.1



An AS May Learn Many Routes

- Multiple connections to neighboring ASes
 - Multiple border routers may learn good routes
 - ... with the same local-pref and AS path length



Hot-Potato (Early-Exit) Routing

- Hot-potato routing
 - Each router selects the closest egress point
 - ... based on the path cost in intradomain protocol
- BGP decision process
 - Highest local preference
 - Shortest AS path
 - **Closest egress point**
 - Arbitrary tie break



← hot potato

